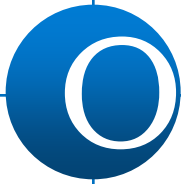


Cycle Problems and Path Problems in Digraphs

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Outline

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- ◆ 02 **Cycle and Path Problems in General**
- ◆ 03 **Degree Conditions and Extremal Digraphs**
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01 Definitions and Notations

Basic Definitions 1

- **Graph G :** an ordered pair $(V(G), E(G))$, in which $V(G)$ is the set of vertices and $E(G)$ is the set of edges, which consists of unordered pairs of vertices.
- **Digraph D :** $(V(D), A(D))$, $V(D)$ set of vertices, $A(D)$ set of arcs, which consists of ordered pairs of vertices.

Basic Definition 2

- **Path.**
- **Cycle.**
- **Length of a path or cycle:** number of edges.
- **Hamiltonian path (cycle):** a path (cycle) that contains every vertex of the graph (digraph).
- **Traceable (Hamiltonian):** a graph (digraph) with a Hamiltonian path (cycle).

Definition & Notations 1

- **Neighbor:** two distinct vertices that are adjacent by an edge (arc) are neighbors.
- **Neighborhood $N(v)$ & Degree $d(v)$:** the set of neighbors of a vertex v , & $d(v) = |N(v)|$.
- **In- & out-neighbor.**
- **In- & out-neighborhood.** $N^-(v), N^+(v)$.
- **In-, out-degree & (total) degree.** $d^-(v) = |N^-(v)|$,
 $d^+(v) = |N^+(v)|$, $d(v) = d^-(v) + d^+(v)$.

Definition & Notations 2

- $\delta(G)$, $\delta(D)$, $\delta^+(D)$ and $\delta^-(D)$: The minimum degree of the vertices in a graph G , minimum (total) degree in a digraph D , the minimum out-degree and in-degree of the vertices in D .
- $\delta^0(D)$: $\min\{\delta^+(D), \delta^-(D)\}$, minimum semi-degree.
- $\sigma_k(G)$: minimum degree sum of all combinations of k independent vertices in a graph G . So $\delta(G) = \sigma_1(G)$.

02

**Cycle and Path
Problems in
General**

Classical Results (Density)

- **(Dirac, 1952)**

If $\delta(G) \geq n/2$, then G is Hamiltonian.

- **(Ore, 1960)**

If $\sigma_2(G) \geq n$, then G is Hamiltonian.

Generalization – Weaker conditions

- **(Fan, 1984)**

If G is 2-connected and for every distinct vertices u and v with $d(u, v) = 2$, $\max\{d(u), d(v)\} \geq n/2$, then G is Hamiltonian.

- **Connectivity:** the minimum number of vertices to be removed from G to make it disconnected.
- **Distance $d(u, v)$:** the length of the shortest path between u and v in G .

Generalization – Various conditions 1

- $\sigma_k(G)$ for $k \geq 3$. $\sigma_k(G) \geq f(n)$.
- **Neighborhood union.**
 $|N(u) \cup N(v)| = d(u) + d(v) - |N(u) \cap N(v)|$.
- **Size:** number of edges. (example. $C(n-1, 2) + 1$ for being Hamiltonian)

Generalization – Various conditions 2

- Degree sequence.
- Forbidden (induced) subgraph.
- Closures.
- Toughness.
- Binding number.

Generalization - More Structures 1

- **Number of Hamiltonian cycles.**
- **Number of edge-disjoint Hamiltonian cycles.**
 - **Decomposition & packing.**
- **Powers of a Hamiltonian cycle.**
- **Long cycles (circumference).**
- **Dominating cycle.**
- **Anti-directed paths and anti-directed cycles.**
- **Factors.**

Generalization - More Structures 2

- **Pancyclicity.**
- **Vertex- (edge-, arc-)pancyclicity.**
- **Hamiltonian-connectedness.**
- **Panconnectedness.**
- **Cycle extendability.**
- **Path extendability.**

03

**Degree Conditions
and Extremal
Digraphs**




Jørgen Bang-Jensen
Gregory Z. Gutin

SPRINGER
MONOGRAPHS IN MATHEMATICS

Digraphs

Theory, Algorithms and Applications
Second Edition



 Springer

Degree Conditions 1

- **(Ghouila-Houri, 1960)**
If $\delta(D) \geq n$, then D is Hamiltonian.
- **(Corollary)**
If $\delta^0(D) \geq n/2$, then D is Hamiltonian.
- **(Woodall, 1972)**
If $d^+(x) + d^-(y) \geq n$ for all pairs of vertices x and y such that there is no arc from x to y , then D is Hamiltonian.

Degree Conditions 2

- **(Meyniel, 1973)**

If D is strong and $d(x) + d(y) \geq 2n - 1$ for all pairs of non-adjacent vertices x and y in D , then D is Hamiltonian.

- **Strong:** for every pair x, y of distinct vertices there exist an (x, y) -path and a (y, x) -path.

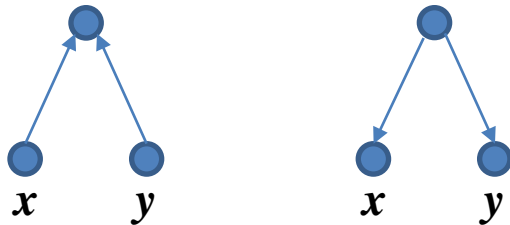
Degree Conditions 3

- **(Bang-Jensen et al. 1996)**

If D is strong and $\min\{d^+(x) + d^-(y), d^-(x) + d^+(y)\} \geq n$

for every pair of dominating non-adjacent and every pair of dominated non-adjacent vertices $\{x, y\}$. Then D is Hamiltonian.

- **Dominating, dominated pair.**



Degree Conditions 4

- **(Conjectured by Bang-Jensen et al. 1996)**

If D is strong and $d(x) + d(y) \geq 2n - 1$ for every pair of dominating non-adjacent and every pair of dominated non-adjacent vertices $\{x, y\}$. Then D is Hamiltonian.

If D is strong and $d(x) + d(y) \geq 2n - 1$ for every pair of dominated non-adjacent vertices $\{x, y\}$. Then D is Hamiltonian.

Extremal Digraphs

- **(Ghouila-Houri, 1960)**
If $\delta(D) \geq n$, then D is Hamiltonian.
- **(Corollary)**
If $\delta^0(D) \geq n/2$, then D is Hamiltonian.
- **(Nash-Williams, 1969)**
Problem: describe all the extreme digraphs for Ghouila-Houri's theorem, i.e. the strong non-Hamiltonian digraphs of order n and minimum degree $n - 1$.

We do not say that the conclusion does not hold. We say that it holds with some exceptions!

– Professor Akira Saito

Extremal Digraphs for Ghouila-Houri's condition

- **(Thomassen, 1981)**

A structural result characterizing strong non-Hamiltonian digraphs of order n and minimum degree $n - 1$.

- **(Darbinyan, 1986)**

If D is a digraph of even order $n \geq 4$ such that $\delta(D) \geq n - 1$ and $\delta^0(D) \geq n/2 - 1$, then either D is Hamiltonian or D belongs to a non-empty finite family of exceptional digraphs.

Extremal Graphs for Woodall's condition

- **(Woodall, 1972)**

if $d^+(x) + d^-(y) \geq n$ for all pairs of vertices x and y such that there is no arc from x to y , then D is Hamiltonian.

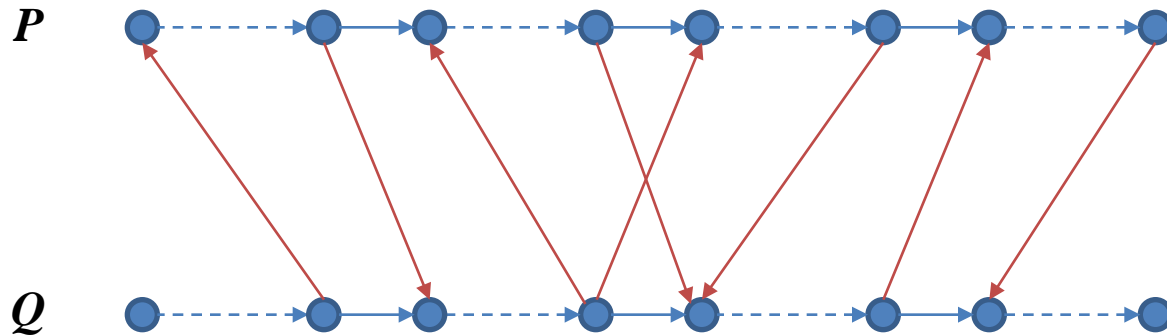
- **(Zhang, Zhang & Wen, SIDMA 2013)**

if $d^+(x) + d^-(y) \geq n - 1$ for all pairs of vertices x and y such that there is no arc from x to y , then D is Hamiltonian, unless D belongs to one of four exceptional families of digraphs.

Techniques

- **Multi-insertion (Bang-Jensen et al., 1996)**

If a path P can be multi-inserted in to a path (cycle, resp.) Q , then there exists a path R with the same starting vertex and ending vertex of Q (a cycle R , resp.), and $V(R)=V(P) \cup V(Q)$.



Further Problems

- **Give a complete characterization of the extremal digraphs of Ghoui-Houri's condition.**
- **Characterize the extremal digraphs of the conditions of Meyniel.**
- **Characterize the extremal digraphs of the conditions of Bang-Jensen et al.**

04 More Structures in Digraphs

The "First Push"

- **(Bondy, 1971)**

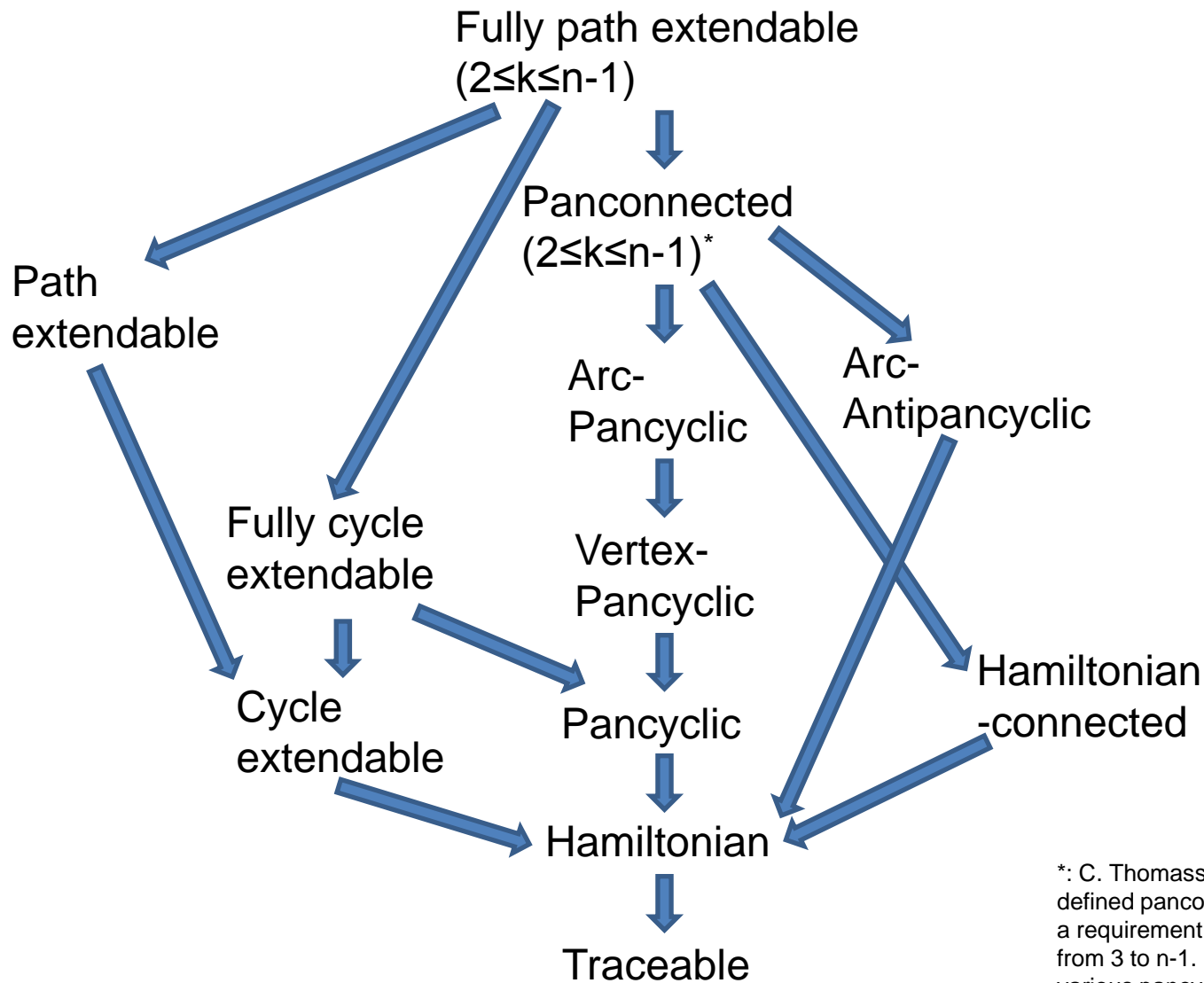
"Meta-conjecture": Almost any nontrivial Hamiltonian condition also implies pancyclicity.

Cycles

- **Pancyclic:** there exist cycles of every length from 3 to n in a graph (digraph).
- **Vertex- (edge-, arc-) pancyclic:** every vertex, (edge, arc) is contained in cycles of every length from 3 to n .
- **Cycle extendable:** for every cycle C , there exists another cycle C' , such that $V(C) \subset V(C')$ and $|V(C')| = |V(C)| + 1$.
- **Fully cycle extendable.**

Paths

- **Hamiltonian-connected:** between every two vertices there is a Hamiltonian path.
- **Panconnected:** between every two vertices there exist paths of every length from $3(2)$ to $n - 1$.
- **Path extendable:** for every path P , there exists another path P' with the same starting and ending vertices, such that $V(P) \subset V(P')$ and $|V(P')| = |V(P)| + 1$.
- **Fully path extendable.**



*: C. Thomassen originally defined panconnecteness with a requirement of path of length from 3 to $n-1$. But to imply various pancyclicities we must have path from 2 to $n-1$ here.

Pancyclicity: Degree Conditions

- **(Thomassen, 1977)**

If $d(x) + d(y) \geq 2n$ whenever x and y are nonadjacent, then either D has cycles of all lengths $2, 3, \dots, n$ or D belongs to some exception classes of digraphs.

- **(Alon & Gutin, 1997)**

If $\delta^0(D) \geq n/2 + 1$, then D is vertex-2-pancyclic. (every vertex is contained in cycles of every length from 2 to n)

- **(Randerath et al., 2002)**

If $\delta^0(D) \geq (n + 1)/2$, then D is vertex-pancyclic.

Pancyclicity: Size Conditions

- **(Häggkvist & Thomassen, 1976)**

Every Hamiltonian digraph on n vertices and at least $n(n + 1)/2 - 1$ arcs is pancyclic.

Cycle Extendability

- **Degree (Hendry, 1989)**

If $\delta(D) \geq 3n/2 - 2$, then D is cycle extendable with some exceptions.

If $\delta^+(D), \delta^-(D) \geq 2n/3 - 1$, then D is cycle extendable with some exceptions.

- **Size (Hendry, 1989)**

If D has at least $(n - 1)^2$ arcs, then D is cycle extendable with some exceptions.

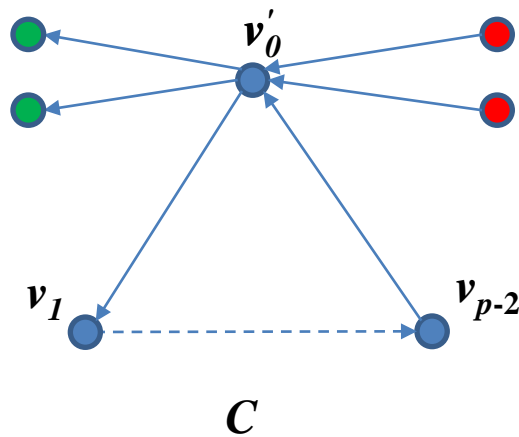
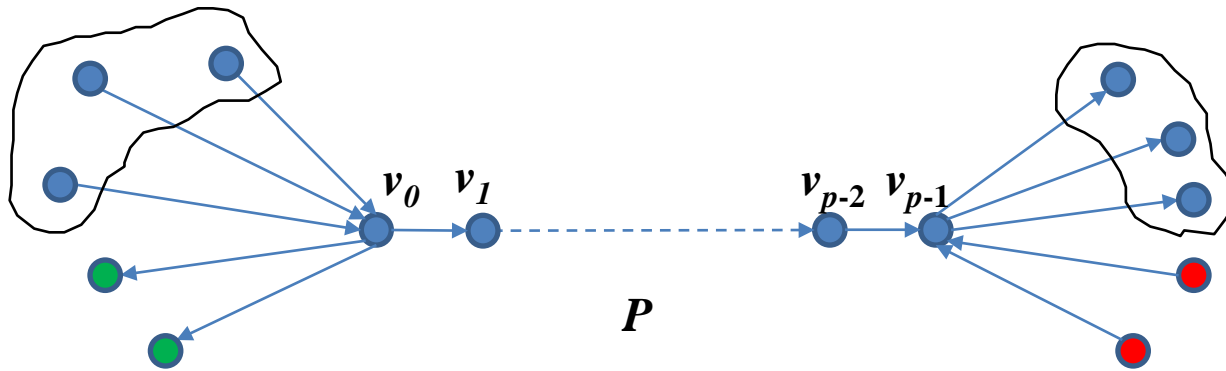
If D is strong and has more than $n^2 - 3n + 4$ arcs, then D is cycle extendable.

Path Extendability

- **Implying relation (Zhang et al., SIDMA 2017)**
If D is path extendable, then D is cycle extendable.
- **Degree (Zhang et al., SIDMA 2017)**
If $\delta(D) \geq 3n/2 - 1$, then D is cycle extendable with some exceptions.
If $\delta^+(D), \delta^-(D) \geq (2n-2)/3$, then D is cycle extendable with some exceptions.
- **Size (Zhang et al., SIDMA 2017)**
If D has at least $(n - 1)^2 + 1$ arcs, then D is path extendable with some exceptions.

Techniques

- **Counting the number of arcs between two parts of the digraph, or counting the degree sums.**
- **A contraction that builds up a corresponding between non-extendable paths and non-extendable cycles.**



P is a non-extendable path in the original digraph, if and only if C is a non-extendable cycle in the resulted digraph.

Further Problems

- **Degree sum conditions for cycle extendability in digraphs.**
- **Degree sum conditions for path extendability in digraphs.**

05 Tournaments and Generalizations

Special Classes of Digraphs

- **Path-mergeable digraphs.**
- **Locally in-semicomplete digraphs.**
- **Oriented graphs: digraphs without 2-cycle.**

Tournaments & Generalizations

- **Tournament:** a digraph with EXACTLY one arc between every two distinct vertices.
- **Bipartite tournament (BT):** a 2-partite tournament.
- **Multipartite tournament (MT, or k -partite tournament):** a digraph whose vertices can be partitioned into k parts, with no arc between vertices in the same part, but exactly one arc between every two vertices in different parts.
- **Semicomplete multipartite digraphs.**

Tournaments 1

- **(Rédei, "Ein kombinatorischer Satz", 1934)**
Every tournament is traceable.
- **(Camion, 1959)**
A strong tournament is Hamiltonian.
- **(Moon, 1968)**
Every strong tournament is vertex-pancyclic.
- **(Moon, 1969 & Hendry, 1989)**
A strong tournament is cycle extendable, unless it belongs to one exceptional class of tournaments.

Tournaments 2

- **(Hendry, 1989)**

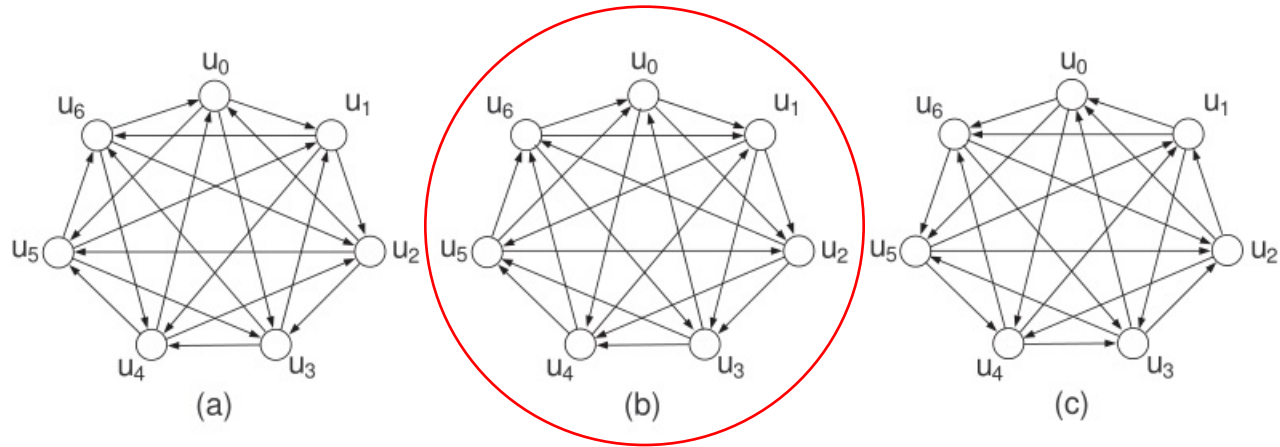
A regular tournament is cycle extendable, unless it belongs to one exceptional class of tournaments.

- **(Zhang et al., SIDMA 2017)**

Tournaments are not generally path extendable.

Neither are strong and regular tournaments.

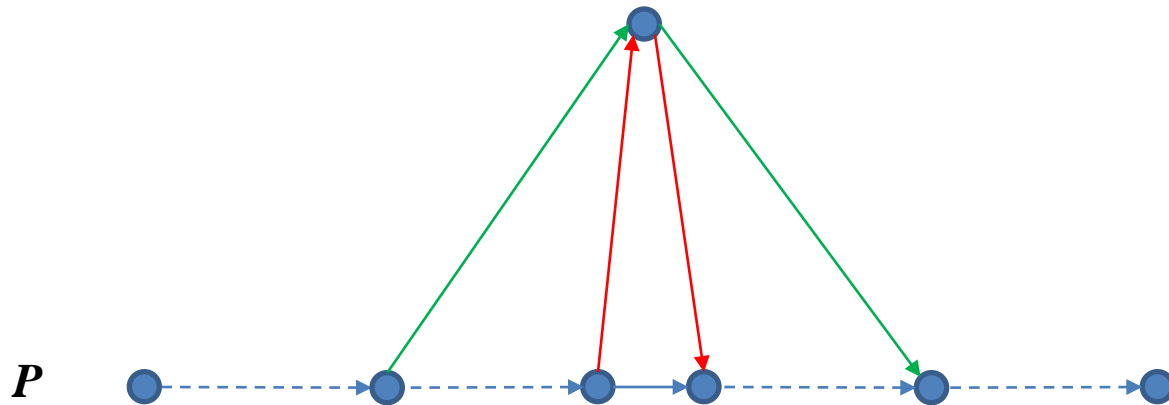
In a regular tournament every path of length at least 2 is extendable, unless it belongs to two exceptional class of tournaments, or isomorphic to a regular tournament on 7 vertices.



The three regular tournaments
on 7 vertices, among which
only (b) is path extendable

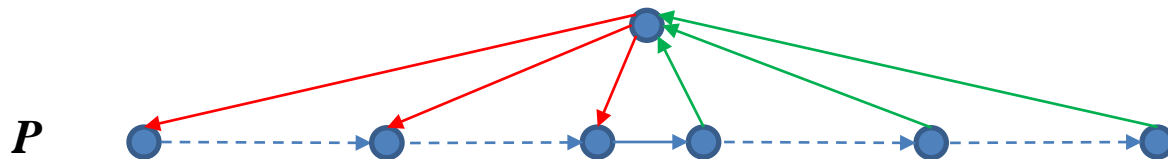
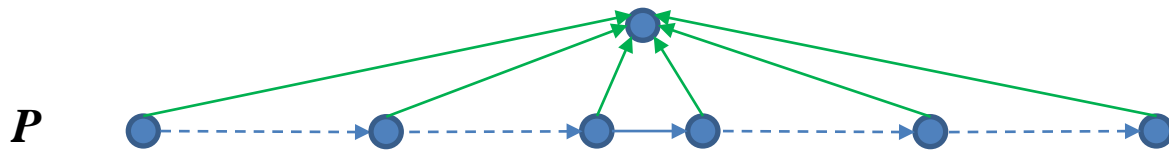
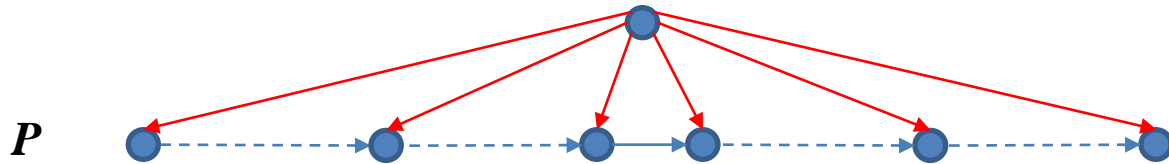
Techniques

- An impossible configure.



Techniques

- Dominating, dominated and hybrid vertices.



Further Problems

- **Generalize some existing conditions for pan-connected in tournaments to those for path extendability in tournaments.**
- **Strong cycle and path properties, such as pancyclicity, cycle extendability, panconnectedness and path extendability of multipartite tournaments.**

**Thanks
for your attention!**